

Computing Circuits and Sign Vectors with Applications to Linear Inequality Systems and Reaction Networks

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Motivation

- $A \in \mathbb{R}^{m \times n}$, $b \in \mathbb{R}^m$

- $x \geq 0$: $x_i \geq 0$ for all i

$$\boxed{\begin{array}{l} Ax = b \\ x \geq 0 \end{array}} \quad x \in \mathbb{R}^n?$$

→ Compute solution (e.g. polyhedra)

- No solution? Certificate? → Dual system

Theorem (Farkas 1902)

Either $\boxed{\begin{array}{l} Ax = b \\ x \geq 0 \end{array}}$ or $\boxed{\begin{array}{l} A^T u \geq 0 \\ b^T u < 0 \end{array}}$ has a solution.

- Theorem of the alternative

$$0 \leq x^T (A^T u) = (Ax)^T u = b^T u < 0 \quad \text{⚡}$$

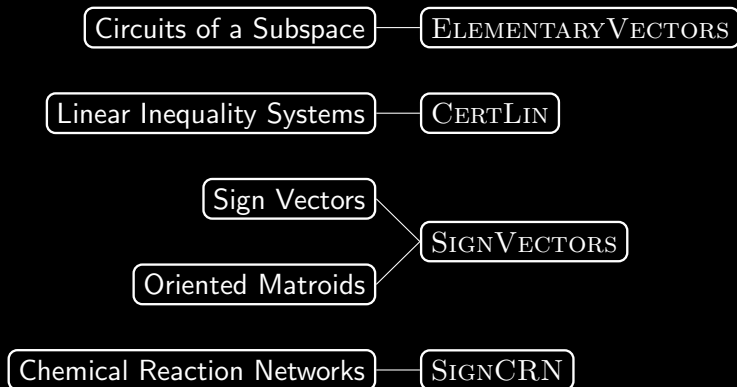
→ u certifies unsolvability!

- Other certificate? Sparse certificate? → Approach via circuits

Approach & Overview

- Results/theory as general as possible
 - General (ordered) fields, integral domains
 - General implementation
 - Exact computation

SAGEMATH Packages



Circuits of a Subspace

- Field $\mathcal{K}$

- For $x \in \mathcal{K}^n$,

$$\text{supp } x = \{i \in [n] : x_i \neq 0\}.$$

$$[n] = \{1, \dots, n\}$$

- Subspace $\mathcal{V}$ of $\mathcal{K}^n$

$$\text{supp} \begin{pmatrix} 1 \\ 0 \\ -2 \end{pmatrix} = \{1, 3\}$$

Definition

$v \in \mathcal{V} \setminus \{0\}$ is **circuit** (of $\mathcal{V}$) if for all $x \in \mathcal{V} \setminus \{0\}$,

$$\text{supp } x \subseteq \text{supp } v \implies \text{supp } x = \text{supp } v.$$

→ **Unique** up to multiples

- Circuits with $\text{supp } x = \text{supp } y \rightarrow x, y$ multiples
- **Finitely many!**

Computing Circuits via Maximal Minors

$$P = \begin{pmatrix} 1 & 2 & 0 & 0 \\ 0 & 1 & 2 & -3 \end{pmatrix} \rightarrow P(\{2, 3\}) = \det \begin{pmatrix} 2 & 0 \\ 1 & 2 \end{pmatrix} = 4$$

- $P = (p_1, \dots, p_n) \in \mathcal{K}^{r \times n}$ with rank r
 - $J = \{j_1, \dots, j_r\} \subseteq [n]$ with $j_1 < \dots < j_r$
- Maximal minor: $P(J) = \det(p_{j_1}, \dots, p_{j_r})$

J	P(J)
{1, 2}	1
{1, 3}	2
{1, 4}	-3
{2, 3}	4
{2, 4}	-6
{3, 4}	0

Theorem

For $I \subseteq [n]$ with $|I| = r + 1$,
define $v \in \mathcal{K}^n$ by

$$v_i = \begin{cases} \varepsilon(i) P(I \setminus \{i\}), & \text{if } i \in I, \\ 0, & \text{else.} \end{cases}$$

If $v \neq 0$, then v is **circuit** of $\ker P$.

- $\varepsilon(i) = \pm 1$

I	{1, 2, 3}	{1, 2, 4}	{1, 3, 4}	{2, 3, 4}
$v(I)$	$\begin{pmatrix} 4 \\ -2 \\ 1 \\ 0 \end{pmatrix}$	$\begin{pmatrix} -6 \\ 3 \\ 0 \\ 1 \end{pmatrix}$	$\begin{pmatrix} 0 \\ 0 \\ 3 \\ 2 \end{pmatrix}$	$\begin{pmatrix} 0 \\ 0 \\ 6 \\ 4 \end{pmatrix}$

- For integral domains!
- **How many?** → $\binom{n}{r+1}$

Unique Enumeration of Circuits

- Reuse zero minor $\rightarrow$ Multiple

- Mark zero minors

- Lexicographic ordering

$\rightarrow$ Algorithm [AR26a]

- Using Grassmann-Plücker identity

Theorem (Grassmann-Plücker Identity)

$K, I \subseteq [n]$ with $|K| = r - 1$ and $|I| = r + 1$. Then,

$$\sum_{i \in I \setminus K} \varepsilon(i) P(K \cup \{i\}) P(I \setminus \{i\}) = 0.$$

- $\varepsilon(i) = \pm 1$

- For commutative rings

Duality of Maximal Minors

- Integral domain $\mathcal{R}$
- $P \in \mathcal{R}^{r \times n}$ with rank r
- $Q \in \mathcal{R}^{n-r \times n}$ with $\ker Q = \text{row } P$

Theorem

There are $\alpha, \beta \in \mathcal{R} \setminus \{0\}$ such that for all $J \subseteq [n]$ with $|J| = r$:

$$\alpha P(J) = \beta \varepsilon(J) Q([n] \setminus J).$$

- $\varepsilon(J) = \pm 1$
- Compute **cocircuits**!
- Circuits of row P

Example

$$P = \begin{pmatrix} 1 & 2 & 0 & 0 \\ 0 & 1 & 2 & -3 \end{pmatrix}$$

$$Q = \begin{pmatrix} 4 & -2 & 1 & 0 \\ -6 & 3 & 0 & 1 \end{pmatrix}$$

J	P(J)	Q([n] \setminus J)
{1, 2}	1	1
{1, 3}	2	-2
{1, 4}	-3	-3
{2, 3}	4	4
{2, 4}	-6	6
{3, 4}	0	0

Package ElementaryVectors

```
sage: from elementary_vectors import *
sage: P = matrix([[1, 2, 0, 0], [0, 1, 2, -3]])
sage: P
[ 1  2  0  0]
[ 0  1  2 -3]
sage: circuits(P)
[(4, -2, 1, 0), (-6, 3, 0, 1), (0, 0, 3, 2)]
sage: cocircuits(P) # circuits in row space
[(0, -1, -2, 3), (1, 0, -4, 6), (2, 4, 0, 0)]

sage: var('a, b')
(a, b)
sage: P = matrix([[1, 2, a, 0], [0, 1, 2, b]])
sage: P
[1 2 a 0]
[0 1 2 b]
sage: circuits(P)
[(-a + 4, -2, 1, 0), (2*b, -b, 0, 1), (a*b, 0, -b, 2),
 (0, a*b, -2*b, -a + 4)]
```

Package ElementaryVectors: Runtime

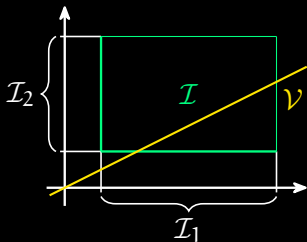
- Maximal minors
 - Compute **once**
 - Hash table
 - No precomputation
 - Compute single circuits

Table: Runtime (in s) for computing circuits

Ring	Size	Circuits	Time	Note
$\mathbb{Z}$	5×15	$5 \cdot 10^3$	0.1	
	5×20	$4 \cdot 10^4$	0.7	
	7×25	$1 \cdot 10^6$	30	4TI2: 180s
	12×25	$5 \cdot 10^6$	300	4TI2: > 2h
$\overline{\mathbb{Q}}$	4×15	$3 \cdot 10^3$	5	
	5×20	$4 \cdot 10^4$	30	only supports

Minty's Lemma

- Subspace $\mathcal{V} \subseteq \mathbb{R}^m$
- Intervals $\mathcal{I} = \mathcal{I}_1 \times \cdots \times \mathcal{I}_m \subseteq \mathbb{R}^m$
- $y \in \mathcal{V} \cap \mathcal{I}$?



Theorem (Rockafellar 1969, Minty 1974)

- (1) Either $y \in \mathcal{V} \cap \mathcal{I}$ exists,
- (2) or **circuit** v of $\mathcal{V}^\perp$ exists with $z^\top v \neq 0$ for all $z \in \mathcal{I}$.

- $\mathbb{R} \rightarrow$ Ordered field $\mathcal{K}$
- **Finitely** many circuits
 - $\rightarrow$ Certify nonexistence of y
- **Linear inequality systems?**
 - $\rightarrow \mathcal{V} = \text{im } M$ with $M \in \mathcal{K}^{m \times n}$
 - $y \in \mathcal{V} \cap \mathcal{I} \rightarrow Mx \in \mathcal{I}$

Example

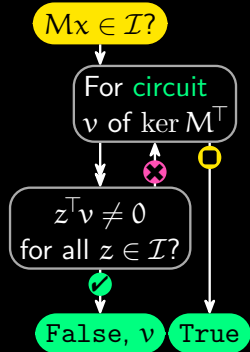
Equivalent are:

- $1 \leq x_1 - x_2 < 2$
 $x_2 > 3$
- $\underbrace{\begin{pmatrix} 1 & -1 \\ 0 & 1 \end{pmatrix}}_M x \in \underbrace{[1, 2) \times (3, \infty)}_{\mathcal{I}}$

Minty's Lemma: Algorithm

Corollary

- (1) Either $Mx \in \mathcal{I}$ has a solution,
 - (2) or **circuit** v of $\ker M^\top$ exists with $z^\top v \neq 0$ for all $z \in \mathcal{I}$.
- Certify **unsolvability**!
 - **Certify solvability**?
 - Apply approach to dual system!
 - Compute solution! **Via circuits**?



□→: End of loop

→→: Parallel

Computing Solutions via Circuits

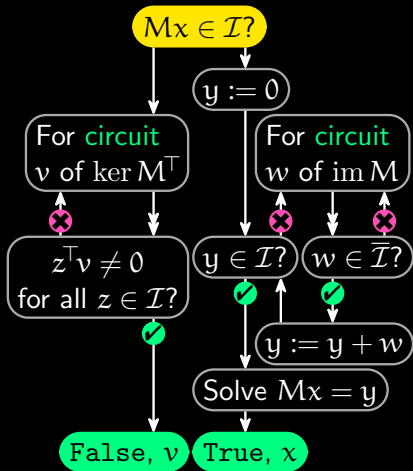
- “Homogeneous” intervals $\mathcal{I}_i \in \{\{0\}, [0, \infty), (0, \infty)\}$
- Closure $\bar{\mathcal{I}}$ with $\bar{\mathcal{I}}_i \in \{\{0\}, [0, \infty)\}$

Corollary (Based on Rockafellar)

If $Mx \in \mathcal{I}$ has a solution,
there are circuits $\mathcal{C}$ of $\text{im } M$ with

$$\sum_{w \in \mathcal{C} \cap \bar{\mathcal{I}}} w = Mx \in \mathcal{I}.$$

- Run algorithms in parallel
- Random circuits
- One hash table for minors
 - For ordered integral domain
 - General systems?
 - Homogenize



CertLin: Live Demo

[Backup]

CertLin: Runtime

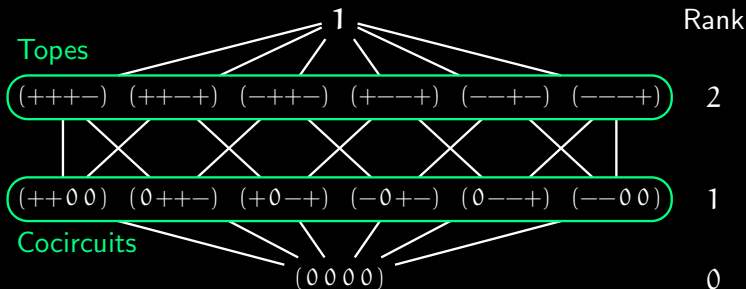
- $Mx \in \mathcal{I} \rightarrow$ Polyhedron?
 - Homogenize
 - Dual system
 - Compute point in polyhedron $\rightarrow$ Certify unsolvability

Table: Runtime (in s) for certifying unsolvability

Ring	Size	Minors	CERTLIN	Polyhedra
$\mathbb{Z}$	20×10	10^5	0.001 – 2	0.7 – 50
	30×15	10^8	0.03 – 3	crash
	60×30	10^{17}	0.5 – 30	–
	100×50	10^{29}	10 – 100	–
$\overline{\mathbb{Q}} \cap \mathbb{R}$	10×5	10^2	0.1 – 0.3	–
	16×8	10^4	0.4 – 4	–

Sign Vectors & Oriented Matroids

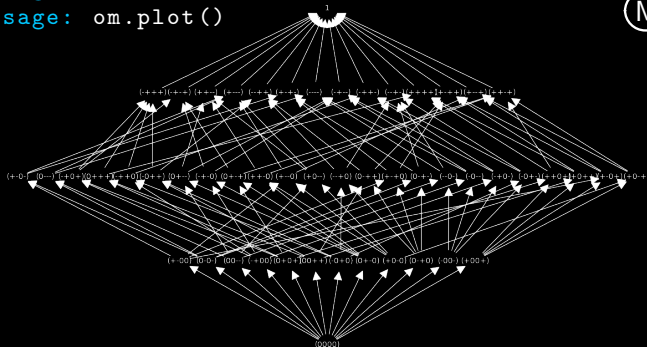
- **Sign vector**: Element in $\{-, 0, +\}^n$ $\text{sign} \begin{pmatrix} -1 \\ 0 \\ 2 \end{pmatrix} = (-0+)$
- **Oriented matroid**: Sign vectors of real subspace
 → Fulfills **covector axioms**
- Partial order on $\{-, 0, +\}^n$ by $0 \preceq -, +$
- Artificial top element **1** → **Big face lattice**



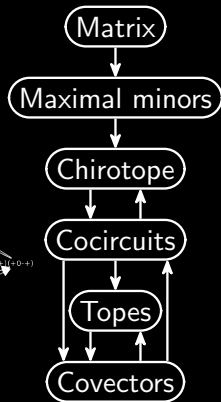
Package SignVectors

- Software presentation [Aic26]
- Fukuda 1991: Topes $\rightarrow$ All covectors by rank
 - Extend: **Construct** big face lattice

```
sage: from sign_vectors import *
sage: P = random_matrix(ZZ, 3, 4)
sage: om = OrientedMatroid(P)
sage: om.plot()
```

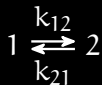


```
sage: om.chirotope() # signs of maximal minors
[+, +, -, 0]
```

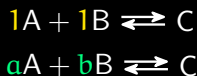


Chemical Reaction Networks

- Weighted directed graph with rate constants $k_{12}, k_{21} > 0$:



- Stoichiometric & kinetic complexes:
 $\mathbf{a}, \mathbf{b} \in \mathbb{R}$.



- ODE system for concentrations $\mathbf{x}$:

$$\frac{d\mathbf{x}}{dt} = \underbrace{\begin{pmatrix} \mathbf{1} & 0 \\ \mathbf{1} & 0 \\ 0 & 1 \end{pmatrix}}_{\mathbf{Y}} \underbrace{\begin{pmatrix} -k_{12} & k_{21} \\ k_{12} & -k_{21} \end{pmatrix}}_{\mathbf{A}_k} \underbrace{\begin{pmatrix} x_A^{\mathbf{a}} x_B^{\mathbf{b}} \\ x_C \end{pmatrix}}_{\mathbf{x}^{\tilde{\mathbf{Y}}}} = \mathbf{Y} \underbrace{\mathbf{A}_k \mathbf{x}^{\tilde{\mathbf{Y}}}}_{=0?}$$

- Positive complex-balanced **equilibrium** (CBE):

$$\mathbf{x} \in \mathbb{R}_{>0}^3 \text{ with } \mathbf{A}_k \mathbf{x}^{\tilde{\mathbf{Y}}} = 0$$

- Study CBEs via subspaces: $\mathcal{S} = \text{im} \begin{pmatrix} -\mathbf{1} \\ -\mathbf{1} \\ 1 \end{pmatrix}, \quad \tilde{\mathcal{S}} = \text{im} \begin{pmatrix} -\mathbf{a} \\ -\mathbf{b} \\ 1 \end{pmatrix}$

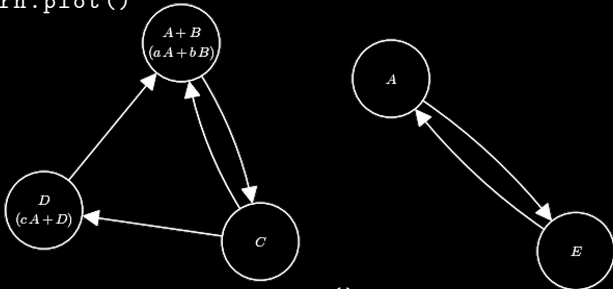
Theorem (Müller, Regensburger 2012)

There is **at most one** positive CBE [...] iff

- $\text{sign}(\mathcal{S}) \cap \text{sign}(\tilde{\mathcal{S}}^\perp) = \{0\}$. $\rightarrow$ Use **chirotope/maximal minors**

Package SignCRN

```
sage: from sign_crn import *
sage: var("a, b, c")
(a, b, c)
sage: species("A, B, C, D, E")
(A, B, C, D, E)
sage: rn = ReactionNetwork()
sage: rn.add_complexes([(0, A+B, a*A+b*B), (1, C)...])
sage: rn.add_reactions([(0, 1), (1, 0), (1, 2)...])
sage: rn.plot()
```



```
sage: rn.has_at_most_one_cbe()
[{a >= 0, b >= 0, a - c >= 0}]
```

Unique Existence of CBE

Theorem (Müller, Hofbauer, Regensburger 2019)

There is *unique* positive CBE [...] iff

- (1) $\text{sign}(\mathcal{S}) \cap \text{sign}(\tilde{\mathcal{S}}^\perp) = \{0\}$,
- (2) for every nonzero $\tilde{\pi} \in \text{sign}(\tilde{\mathcal{S}}^\perp)_\oplus$,
there is nonzero $\pi \in \text{sign}(\mathcal{S}^\perp)_\oplus$ with $\pi \preceq \tilde{\pi}$, $\rightarrow$ Use cocircuits
- (3) and $(\mathcal{S}, \tilde{\mathcal{S}})$ is nondegenerate. $\rightarrow$???

$\rightarrow$ Recursive algorithm for (3) [AMR24]

- Nonnegative cocircuits
- Linear inequality systems

Summary & Outlook

- Packages:

<https://github.com/MarcusAichmayr>

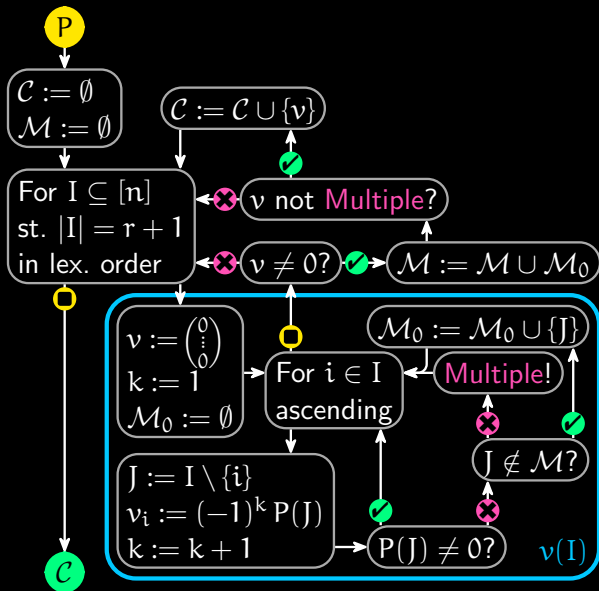


- C++ libraries
- CERTLIN paper in preparation [AR26b]
- Implement more sign vector conditions for CRNs
- Study conditions for positive solutions of polynomial equations

References

- [Aic26] Marcus S. Aichmayr. “SignVectors: a SageMath Package for Oriented Matroids”. *Accepted. Software presentation ISSAC. 2026.*
- [AMR24] Marcus S. Aichmayr, Stefan Müller, and Georg Regensburger. “A SageMath package for elementary and sign vectors with applications to chemical reaction networks”. In: *Mathematical software—ICMS 2024*. Vol. 14749. LNCS. Springer, Cham, 2024, pp. 155–164.
- [AR26a] Marcus S. Aichmayr and Georg Regensburger. “Computing Circuits over Integral Domains”. *Submitted. 2026.*
- [AR26b] Marcus S. Aichmayr and Georg Regensburger. “How to Certify Solvability of Linear Inequality Systems with Circuits”. *In preparation. 2026.*

Unique Enumeration of Circuits



Homogeneous Linear Inequality Systems

Theorem (Motzkin 1936)

Exactly one of the following has a solution:

$$(H1) \quad \begin{cases} Ax > 0 \\ Bx \geq 0 \\ Cx = 0 \end{cases}$$

$$(H2) \quad \begin{cases} A^T u + B^T v + C^T w = 0 \\ u, v \geq 0, \quad u \neq 0 \end{cases}$$

$\rightarrow Mx \in \mathcal{I}$

$$(H1') \quad \begin{pmatrix} A \\ B \\ C \end{pmatrix} x \in (0, \infty)^p \times [0, \infty)^q \times \{0\}^r$$

$$(H2') \quad \begin{pmatrix} A^T & B^T & C^T \\ \text{Id} & 0 & 0 \\ 0 & \text{Id} & 0 \\ 1^T & 0 & 0 \end{pmatrix} \begin{pmatrix} u \\ v \\ w \end{pmatrix} \in \{0\}^n \times [0, \infty)^{p+q} \times (0, \infty)$$

Homogenization of Linear Inequality Systems

Lemma (Homogenization)

The following are equivalent:

$$(1) \quad \begin{cases} Ax < a \\ Bx \leq b \end{cases}$$

$$(2) \quad \begin{cases} \begin{pmatrix} A & a \\ 0 & 1 \end{pmatrix} \begin{pmatrix} x \\ \tau \end{pmatrix} > 0 \\ \begin{pmatrix} B & b \end{pmatrix} \begin{pmatrix} x \\ \tau \end{pmatrix} \geq 0 \end{cases}$$

Proof.

If x is a solution of (1), then $(-x, 1)^\top$ is a solution of (2).

If $(x, \tau)^\top$ is a solution of (2), then $-\frac{x}{\tau}$ is a solution of (1).



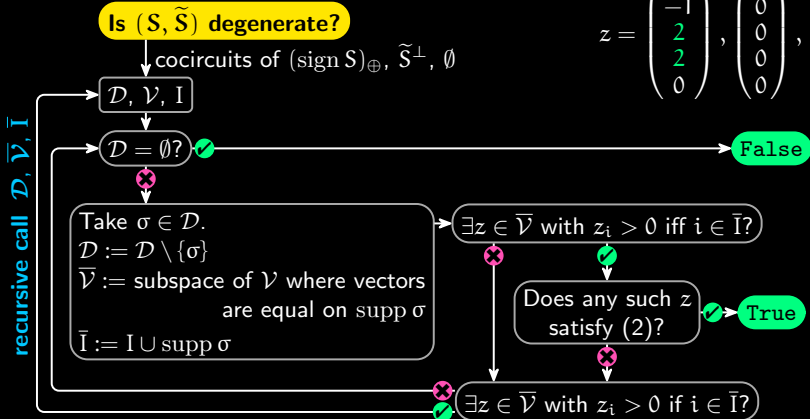
Exactly One CBE – Recursive Algorithm

Definition

$(\mathcal{S}, \tilde{\mathcal{S}})$ is **degenerate** if there is $z \in \tilde{\mathcal{S}}^\perp$:

- (1) There is nonempty set of cocircuits $\mathcal{D} \subseteq \text{sign}(\mathcal{S})_\oplus$:
 - For all $\sigma \in \mathcal{D}$ and for all $i, j \in \text{supp } \sigma$, $z_i = z_j > 0$.
 - For all $i \notin \bigcup_{\sigma \in \mathcal{D}} \text{supp } \sigma$, $z_i \leq 0$.
- (2) For every cocircuits $\sigma \in \text{sign}(\mathcal{S}^\perp)_\oplus$, $\text{supp } \sigma \not\subseteq \text{supp } z$.

$$z = \begin{pmatrix} 1 \\ -1 \\ 2 \\ 2 \\ 0 \end{pmatrix}, \quad \overbrace{\begin{pmatrix} + \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}}^{\in \mathcal{D}}, \quad \begin{pmatrix} 0 \\ 0 \\ + \\ + \\ 0 \end{pmatrix}$$



Package CertLin

```
sage: from certlin import *
sage: M = random_matrix(ZZ, 15, 5)
sage: I = Intervals.random(15, ring=ZZ)
sage: S = LinearInequalitySystem(M, I)
sage: S
[ 0  5  1 -1  1] x in [1, 2)
[-1  1  0  1  0] x in [0, 2]
[ 0  2  3 -3  0] x in [1, 3]
[-1 -5  0  0  1] x in (-1, 0]
[-1 -5 -1  0  3] x in (-oo, -1)
[ 0  1 -5 -1  1] x in [1, +oo)
[-2  1  1  1  0] x in (-1, 0]
[ 2  1 -1 -1  1] x in (-oo, +oo)
[-1  1  0 -4 -1] x in (-2, -1)
[ 1  0 -2  1  1] x in (0, 1]
[ 4 -3  3 -3  0] x in [-1, 1)
[-1  1  0  2  1] x in (-1, +oo)
[ 1  0 -2 -2  1] x in (-oo, 1)
[-1 -2 -4  0 -3] x in {0}
[ 1 -4 -2 -1  0] x in (-oo, -2]
```

CertLin: Certify

```
sage: S.certify_unsolvability() # sparse
(134, 0, 0, 356, -216, 14, 0, 0, 0, 144, 0, 0, 0, 0, -4)

sage: S.certify() # both algorithms
(False, (134, 0, 0, 356, -216, 14, 0, 0, 0, 144, 0, 0, 0, 0, -4))

sage: S.certify(random=True) # faster
(False, (0, -1, 9, 0, 0, 25, -12, 0, -30, -55, 0, 0, 0, 0, 0))
```

CertLin: Homogenization

```
sage: S.to_homogeneous()
[ 0  5  1 -1  1 -2] x > 0
[ 1  5  0  0 -1 -1] x > 0
[-1 -5 -1  0  3  1] x > 0
[ 2 -1 -1 -1  0 -1] x > 0
[ 1 -1  0  4  1 -2] x > 0
[-1  1  0 -4 -1  1] x > 0
[-1  0  2 -1 -1  0] x > 0
[ 4 -3  3 -3  0 -1] x > 0
[ 1 -1  0 -2 -1 -1] x > 0
[ 1  0 -2 -2  1 -1] x > 0
[ 0  0  0  0  0 -1] x > 0
[ 0 -5 -1  1 -1  1] x >= 0
[ 1 -1  0 -1  0  0] x >= 0
[-1  1  0  1  0 -2] x >= 0
[ 0 -2 -3  3  0  1] x >= 0
[ 0  2  3 -3  0 -3] x >= 0
[-1 -5  0  0  1  0] x >= 0
[ 0 -1  5  1 -1  1] x >= 0
[-2  1  1  1  0  0] x >= 0
[ 1  0 -2  1  1 -1] x >= 0
[-4  3 -3  3  0 -1] x >= 0
[ 1 -4 -2 -1  0  2] x >= 0
[-1 -2 -4  0 -3  0] x = 0
```

CertLin: Dual System

sage: S.dual()

```
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  1  1  1  1  1  1  1  1  1  1] x > 0
[ 1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0  0] x >= 0
[ 0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  1  0  0  0  0] x >= 0
[ 0  1 -1  0  0 -1  0 -2  1 -4  1 -1  1  0  1 -1  2  1 -1 -1  4  1  1  0] x = 0
[-5 -1  1 -2  2 -5 -1  1  0  3  2 -2 -4  5  5 -5 -1 -1  1  0 -3 -1  0  0] x = 0
[-1  0  0 -3  3  0  5  1 -2 -3  4 -4 -2  1  0 -1 -1  0  0  2  3  0 -2  0] x = 0
[ 1 -1  1  3 -3  0  1  1  1  3  0  0 -1 -1  0  0 -1  4 -4 -1 -3 -2 -2  0] x = 0
[-1  0  0  0  0  1 -1  0  1  0  3 -3  0  1 -1  3  0  1 -1 -1  0 -1  1  0] x = 0
[ 1  0 -2  1 -3  0  1  0 -1 -1  0  0  2 -2 -1  1 -1 -2  1  0 -1 -1 -1 -1] x = 0
```

CertLin: Real Algebraic Numbers

```
sage: M = matrix([[1, sqrt(3), -1/2, 0], ...])
sage: I = Intervals.from_bounds([1, -oo, ...], ...)
sage: S = LinearInequalitySystem(M, I)
sage: S
[      1  sqrt(3)      -1/2      0]  x in [1, sqrt(5))
[      1      0 -sqrt(5)      3]  x in (-oo, 0)
[ 3^(1/3)      0      5/2      1]  x in [sqrt(2), +oo)
```

```
sage: S.certify_unsolvability() # instant
```

```
...
```

```
ValueError: A solution exists.
```

```
sage: S.certify()
```

```
(True, (0, 1/30*(12*sqrt(5)*3^(5/6)*sqrt(2) + 58*sqrt(5)*3^(5/6) + 85*sqrt(5)*
sqrt(3) + 30*sqrt(3)*sqrt(2) + 170*3^(5/6) + 145*sqrt(3))/(6*sqrt(5)
*3^(1/3) + 2*sqrt(5) + 4*3^(1/3) + 15), 2/5*(6*sqrt(5)*3^(1/3)*sqrt(2) + 4*
sqrt(5)*3^(1/3) + 15*sqrt(2) + 10)/(6*sqrt(5)*3^(1/3) + 2*sqrt(5) +
4*3^(1/3) + 15), 1/2*(4*sqrt(5)*sqrt(2) - 2*sqrt(5)*3^(1/3) + 8*3^(1/3)*
sqrt(2) + 2*sqrt(5) + 4*3^(1/3) - 5)/(6*sqrt(5)*3^(1/3) + 2*sqrt(5) +
4*3^(1/3) + 15)))
```

```
sage: S.find_solution()
```

```
(0, 1/30*(12*sqrt(5)*3^(5/6)*sqrt(2) + 58*sqrt(5)*3^(5/6) + 85*sqrt(5)...
```